Program Representations

Representing programs

- Goals

Representing programs

- Primary goals
  - analysis is easy and effective
    - just a few cases to handle
    - directly link related things
  - transformations are easy to perform
  - general, across input languages and target machines
- Additional goals
  - compact in memory
  - easy to translate to and from
  - tracks info from source through to binary, for source-level debugging, profilling, typed binaries
  - extensible (new opts, targets, language features)
  - displayable

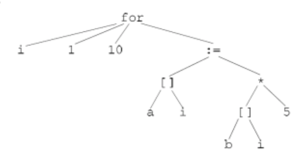
Option 1: high-level syntax based IR

- Represent source-level structures and expressions directly
- Example: Abstract Syntax Tree

Source:

```
for i := 1 to 10 do
  a[i] := b[i] * 5;
end
```

AST:



Option 2: low-level IR

- Translate input programs into low-level primitive chunks, often close to the target machine
- Examples: assembly code, virtual machine code (e.g. stack machines), three-address code, register-transfer language (RTL)

- Standard RTL instrs:

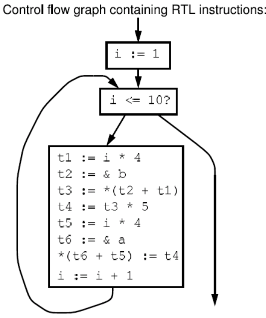
assignment	`x := y;`
unary op	`x := op y;`
binary op	`x := y op z;`
address-of	`p := &y;`
load	`x := *(p + o);`
store	`*(p + o) := x;`
call	`x := f(...);`
unary compare	`op x ?`
binary compare	`x op y ?`

Option 2: low-level IR

Source:

```
for i := 1 to 10 do
  a[i] := b[i] * 5;
end
```

Control flow graph containing RTL instructions:



Comparison

Comparison

- Advantages of high-level rep
  - analysis can exploit high-level knowledge of constructs
  - easy to map to source code (debugging, profiling)
- Advantages of low-level rep
  - can do low-level, machine specific reasoning
  - can be language-independent
- Can mix multiple reps in the same compiler

Components of representation

- Control dependencies: sequencing of operations
  - evaluation of if & then
  - side-effects of statements occur in right order
- Data dependencies: flow of definitions from defs to uses
  - operands computed before operations
- Ideal: represent just dependencies that matter
  - dependencies constrain transformations
  - fewest dependences ⇒ flexibility in implementation

1 a := b + c

2 d := e + f

Control dependencies

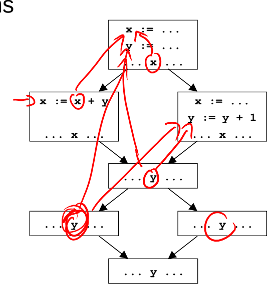
- Option 1: high-level representation
  - control implicit in semantics of AST nodes
- Option 2: control flow graph (CFG)
  - nodes are individual instructions
  - edges represent control flow between instructions
- Options 2b: CFG with basic blocks
  - basic block: sequence of instructions that don't have any branches, and that have a single entry point
  - BB can make analysis more efficient: compute flow functions for an entire BB before start of analysis

Control dependencies

- CFG does not capture loops very well
- Some fancier options include:
  - the Control Dependence Graph
  - the Program Dependence Graph
- More on this later. Let's first look at data dependencies

Data dependencies

- Simplest way to represent data dependencies: def/use chains

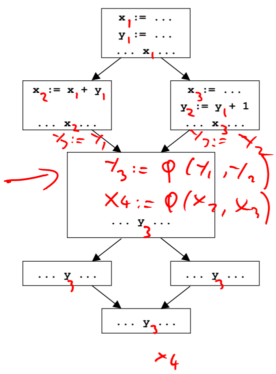


Def/use chains

- Directly captures dataflow
  - works well for things like constant prop
- But...
- Ignores control flow
  - misses some opt opportunities since conservatively considers all paths
  - not executable by itself (for example, need to keep CFG around)
  - not appropriate for code motion transformations
- Must update after each transformation
- Space consuming

SSA

- Static Single Assignment
  - invariant: each use of a variable has only one def



SSA

- Create a new variable for each def
- Insert ϕ pseudo-assignments at merge points
- Adjust uses to refer to appropriate new names

- Question: how can one figure out where to insert ϕ nodes using a liveness analysis and a reaching defns analysis.

Converting back from SSA

- Semantics of $x_3 := \phi(x_1, x_2)$
  - set x_3 to x_i if execution came from ith predecessor
- How to implement ϕ nodes?

Converting back from SSA

- Semantics of $x_3 := \phi(x_1, x_2)$
  - set x_3 to x_i if execution came from ith predecessor
- How to implement ϕ nodes?
  - Insert assignment $x_3 := x_1$ along 1st predecessor
  - Insert assignment $x_3 := x_2$ along 2nd predecessor
- If register allocator assigns x_1, x_2 and x_3 to the same register, these moves can be removed
  - x_1 .. x_n usually have non-overlapping lifetimes, so this kind of register assignment is legal

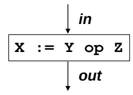
Recall: Common Sub-expression Elim

- Want to compute when an expression is available in a var
- Domain: $\{x \to E_1,\ y \to E_2,\ z \to E_3\}$

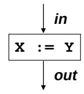
$S = \{x \to E \mid x \in Var,\ E \in Expr\}$

$D = 2^S$
$\perp = S$
$T = \emptyset$
$\sqcup = \cap$

Recall: CSE Flow functions

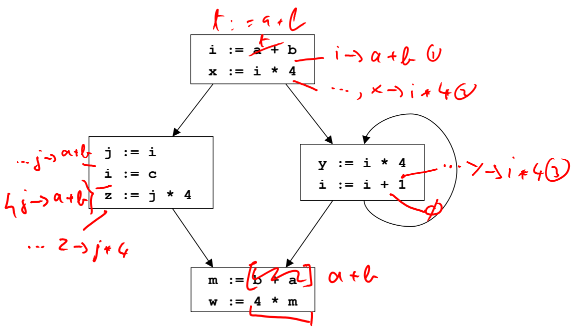


$F_{X := Y \text{ op } Z}(in) = in - \{ X \to * \}$
$- \{ * \to ... X ... \} \cup$
$\{ X \to Y \text{ op } Z \mid X \neq Y \wedge X \neq Z\}$



$F_{X := Y}(in) = in - \{ X \to * \}$
$- \{ * \to ... X ... \} \cup$
$\{ X \to E \mid Y \to E \in in \}$

Example



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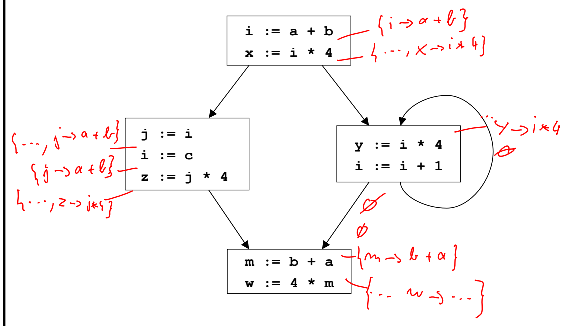

Problems

- z := j * 4 is not optimized to z := x, even though x contains the value j * 4
- m := b + a is not optimized, even though a + b was already computed
- w := 4 * m it not optimized to w := x, even though x contains the value 4 *m

Problems: more abstractly

- Available expressions overly sensitive to name choices, operand orderings, renamings, assignments
- Use SSA: distinct values have distinct names
- Do copy prop before running available exprs
- Adopt canonical form for commutative ops

Example in SSA

in

`X := Y op Z`

out

$F_{X := Y\ op\ Z}(in) = in \cup \{ X \to Y_{op} Z \}$

in_0 in_1

`X := ϕ(Y,Z)`

out

$F_{X := \phi(Y,Z)}(in_0, in_1) =$

$a_1 := E$ $a_2 := E$

$a_3 := \phi(a_1, a_2)$

Example in SSA

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`X := Y op Z`

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$F_{X := Y\ op\ Z}(in) = in \cup \{ X \to Y \text{ op } Z \}$

in_0 in_1

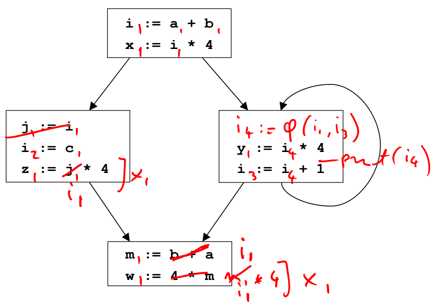
`X := ϕ(Y,Z)`

out

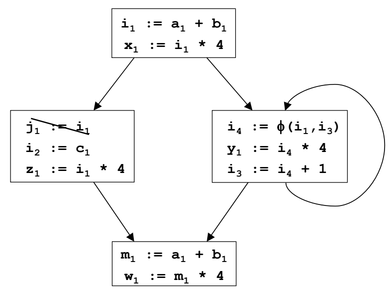
$F_{X := \phi(Y,Z)}(in_0, in_1) = (in_0 \cap in_1) \cup$

$\{ X \to E \mid Y \to E \in in_0 \wedge Z \to E \in in_1 \}$

Example in SSA



Example in SSA



What about pointers?

- Pointers complicate SSA. Several options.
- Option 1: don’t use SSA for pointed to variables
- Option 2: adapt SSA to account for pointers
- Option 3: define src language so that variables cannot be pointed to (eg: Java)

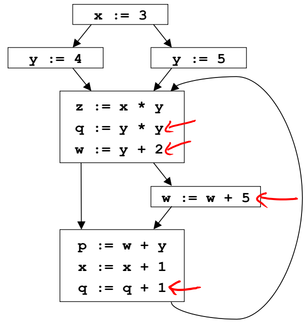
SSA helps us with CSE

- Let’s see what else SSA can help us with
- Loop-invariant code motion

Loop-invariant code motion

- Two steps: analysis and transformations
- Step1: find invariant computations in loop
  - invariant: computes same result each time evaluated
- Step 2: move them outside loop
  - to top if used within loop: **code hoisting**
  - to bottom if used after loop: **code sinking**

Example



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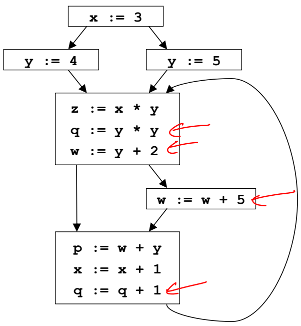


Detecting loop invariants

- An expression is invariant in a loop L iff:

(base cases)

2 and x

- it's a constant
- it's a variable use, all of whose defs are outside of L

$a = b \Rightarrow f(a) = f(b)$

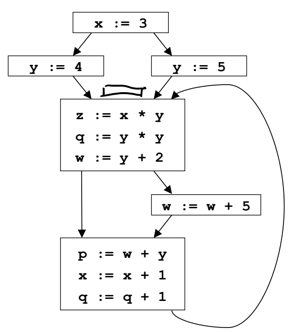
(inductive cases)

- it's a pure computation all of whose args are loop-invariant
- it's a variable use with only one reaching def, and the rhs of that def is loop-invariant

Computing loop invariants

- Option 1: iterative dataflow analysis
  - optimistically assume all expressions loop-invariant, and propagate
- Option 2: build def/use chains
  - follow chains to identify and propagate invariant expressions
- Option 3: SSA
  - like option 2, but using SSA instead of def/use chains

Example using def/use chains



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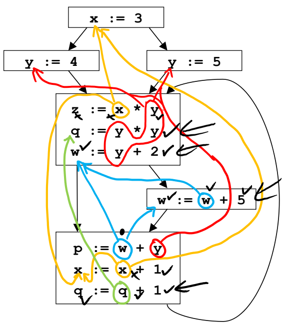

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Loop invariant detection using SSA

- An expression is invariant in a loop L iff:

(base cases)

- it's a constant
- it's a variable use, all of whose **single** defs are outside of L

(inductive cases)

- it's a pure computation all of whose args are loop-invariant
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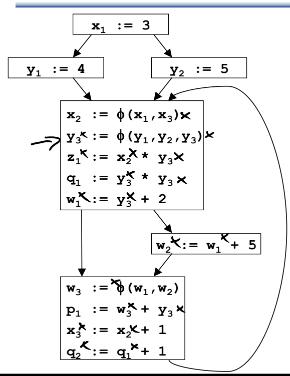
- ϕ functions are not pure

Example using SSA



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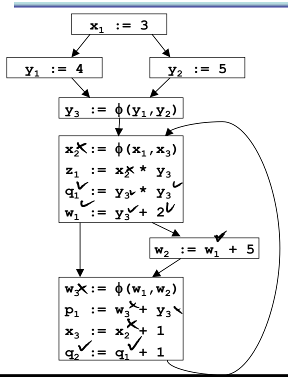
Example using SSA and preheader



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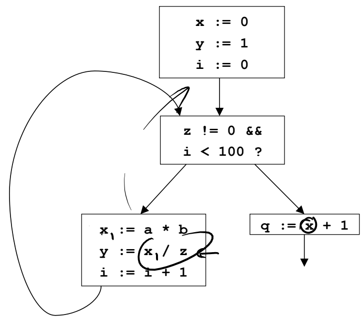
Summary: Loop-invariant code motion

- Two steps: analysis and transformations
- Step1: find invariant computations in loop
  - invariant: computes same result each time evaluated
- Step 2: move them outside loop
  - to top if used within loop: **code hoisting**
  - to bottom if used after loop: **code sinking**

Code motion

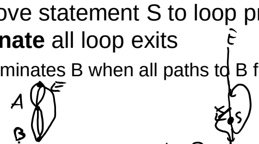
- Say we found an invariant computation, and we want to move it out of the loop (to loop pre-header)
- When is it legal?
- Need to preserve relative order of invariant computations to preserve data flow among move statements
- Need to preserve relative order between invariant computations and other computations

Example



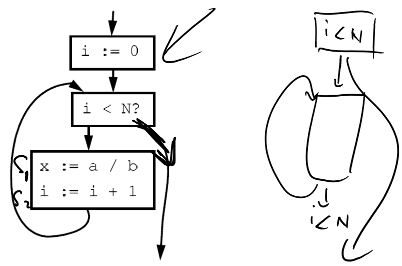
Lesson from example: domination restriction

- To move statement S to loop pre-header, S must **dominate** all loop exits
  [A dominates B when all paths to B first pass through A]

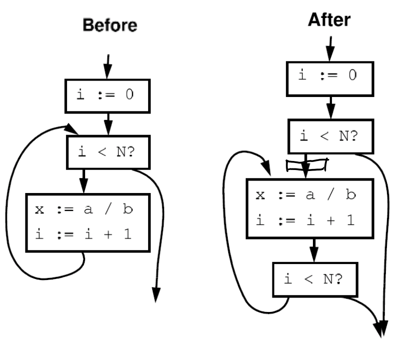


- Otherwise may execute S when never executed otherwise
- If S is pure, then can relax this constraint at cost of possibly slowing down the program

Domination restriction in for loops



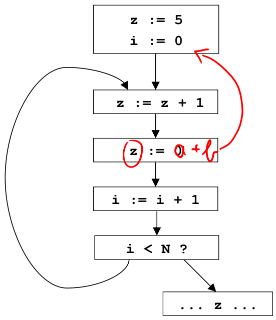
Domination restriction in for loops



Avoiding domination restriction

- Domination restriction strict
  - Nothing inside branch can be moved
  - Nothing after a loop exit can be moved
- Can be circumvented through loop normalization
  - while-do => if-do-while

Another example



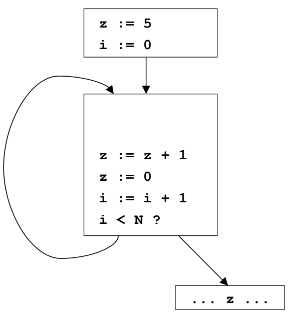
Data dependence restriction

- To move S: **z := x op y**:

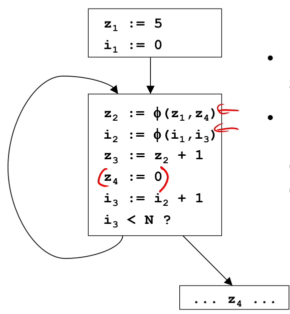
  S must be the only assignment to **z** in loop, and no use of **z** in loop reached by any def other than S

- Otherwise may reorder defs/uses

Avoiding data restriction



Avoiding data restriction



$z_4 := 0$

... z_4 ...

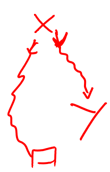
- Restriction much simpler in SSA!!!
- Implementation of phi nodes as moves will cope with re-ordered defs/uses

Summary of Data dependencies

- We've seen SSA, a way to encode data dependencies better than just def/use chains
  - makes CSE easier
  - makes loop invariant detection easier
  - makes code motion easier
- Now we move on to looking at how to encode control dependencies

Control Dependencies

- A node (basic block) Y is control-dependent on another X iff X determines whether Y executes
  - there exists a path from X to Y s.t. every node in the path other than X and Y is post-dominated by Y
  - X is not post-dominated by Y



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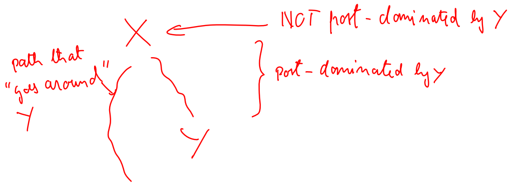
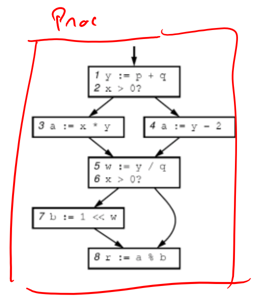
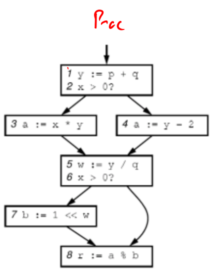


Example

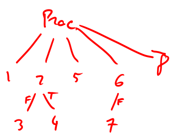


4 CD on 2
3 " " "
7 " " 6
1,2,5,6,8 CD

Example



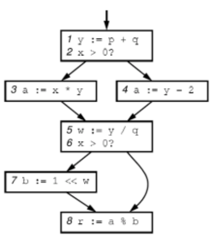
Control dependence relation
3 depends on 2
4 " " 2
7 " " 6



Control Dependence Graph

- Control dependence graph: Y descendent of X iff Y is control dependent on X
  - label each child edge with required condition
  - group all children with same condition under region node
- Program dependence graph: super-impose dataflow graph (in SSA form or not) on top of the control dependence graph

Example



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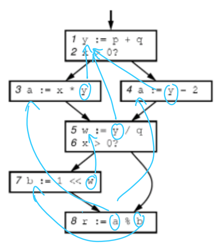

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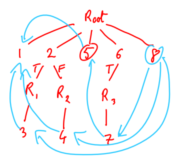
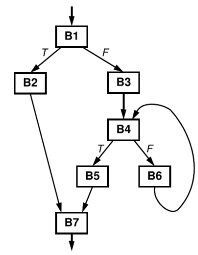


Another example



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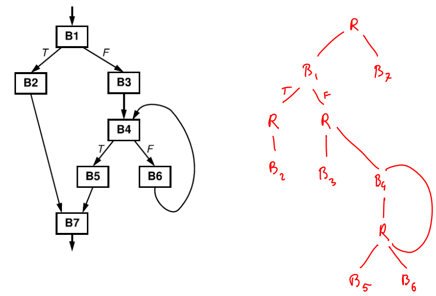
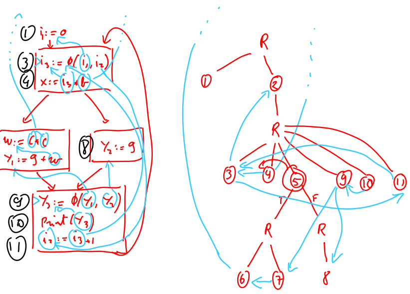


Another example

```
1  i1 := 0;
   while 2 ... do
3    i3 := φ(i1, i2);
4    x := i3 * b;
     if 5 ... then
6      w := c * c;
7      y1 := 9 + w;
     else
8      y2 := 9;
     end
9    y3 := φ(y1, y2);
10   print(y3);
11   i2 := i3 + 1;
   end
```



Summary of Control Depence Graph

- More flexible way of representing control-depencies than CFG (less constraining)
- Makes code motion a local transformation
- However, much harder to convert back to an executable form

Course summary so far

- Dataflow analysis
  - flow functions, lattice theoretic framework, optimistic iterative analysis, precision, MOP
- Advanced Program Representations
  - SSA, CDG, PDG
- Along the way, several analyses and opts
  - reaching defns, const prop & folding, available exprs & CSE, liveness & DAE, loop invariant code motion
- Next: dealing with pointers