# CSE130/230 - WEEK 5 DI 

Interpreters, PA4, and beyond

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## The Plan for Today

1. Interpreters
2. PA 4 Overview
3. PA4 Concepts

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1. Interpreters
2. PA 4 Overview
3. PA4 Concepts
a. Environments
b. Closures
c. Apps
i. Let, Letrec
d. Native ops

Obligatory Halloween meme

## Interpreters

~Of the realms beyond~

## What is an Interpreter?

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The expression you're currently evaluating

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Usually, it consists of an evaluation loop that recursively resolves the arguments to an operator from expressions to values.

## let rec eval (evn,e) = failwith "to be written"



The environment: array of tuples of the form ("var", "value")

## How to implement an interpreter?

let rec eval (evn,e) = failwith "to be written"

## How to implement an interpreter?

## let rec eval (evn,e) =

Pattern match `e` with the data constructors and handle each case;;
Sometimes add a new variable to `env`

Also check that types are correct: cannot do 4 + "Burger", for example

## The Big Picture



## The Big Picture



# Environments 

~Of the undead~

## Let's run some code in our heads!

let $\mathrm{a}=1 \mathrm{in}$
let $b=2$ in
let $\mathrm{a}=\mathrm{a}+1 \mathrm{in}$
$a+b$

## Let's run some code in our heads!

let $\mathrm{a}=1 \mathrm{in}$
let $b=2$ in
let $\mathrm{a}=\mathrm{a}+1 \mathrm{in}$
$a+b$
What's the value of the final expression

## Let's run some code in our heads!

let $\mathbf{a}=1$ in

$$
\begin{aligned}
& \text { let } b=2 \text { in } \\
& \qquad \text { let } a=\mathbf{a}+1 \text { in }
\end{aligned}
$$

$$
a+b / / 4
$$

## Let's run some code in our heads!

let $\mathbf{a}=1$ in

$$
\begin{aligned}
& \text { let } b=2 \text { in } \\
& \qquad \begin{array}{r}
\text { let } a=a+1 \text { in } \\
a+b / / 4
\end{array}
\end{aligned}
$$



## How is the environment filled?

let $\mathbf{a}=1$ in

$$
\begin{aligned}
& \text { let } b=2 \text { in } \\
& \qquad \text { let } \mathrm{a}=\mathrm{a}+1 \text { in }
\end{aligned}
$$

$$
\text { + b // } 4
$$

## How is the environment filled?

let $\mathbf{a}=1$ in

$$
\begin{aligned}
& \text { let } b=2 \text { in } \\
& \qquad \begin{array}{l}
\text { let } a=a+1 \\
\text { in } \\
a+b / / 4
\end{array}
\end{aligned}
$$

## How is the environment filled?

let $\mathbf{a}=1$ in

```
let b = 2 in
    let a = a + 1 in
```

    + b // 4
    
## How is the environment filled?

let $\mathbf{a}=1$ in

$$
\begin{aligned}
& \text { let } \mathrm{b}=\mathbf{2} \text { in } \\
& \qquad \text { let } \mathrm{a}=\mathbf{a}+\mathbf{1} \text { in }
\end{aligned}
$$

$$
\text { + b // } 4
$$

## How is the environment filled?

$$
\text { let } \mathrm{a}=1 \text { in }
$$

$$
\text { let } \mathrm{b}=\mathbf{2} \text { in }
$$

$$
\text { let } \mathrm{a}=\mathbf{a}+1 \text { in }
$$

$$
+b / / 4
$$

## Environment

ListAssoc finds the left-most definition of any variable in the environment. So will resolve to $2+2=4$ instead of $1+2=3$

## Closures

## From the beyond

## Closures

Construction: Closure(env, name, argument, body)

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Either: None or Some 'name'

## Closures

Construction: Closure(env, name, argument, body)
This is the representation of a function in your environment

## Either: None or Some 'name'

## When is a closure created?

let $\mathrm{a}=1 \mathrm{in}$
let $b=2$ in
let foo $=$ fun $x->x+1$

$$
\text { in } x+b+a+1
$$

## When is a closure created?

```
let a = 1 in
let b = 2 in
    let foo = fun x -> x + 1
        in x + b + a + 1
```


## When is a closure created?

let $\mathbf{b}=\mathbf{2}$ in
let foo = fun $\mathbf{x}$-> $\mathbf{x}+1$
in $\mathbf{x}+\mathrm{b}+\mathrm{a}+1$

## When is a closure created?

let $\mathbf{b}=\mathbf{2}$ in
let foo $=$ fun $\mathbf{x}$-> $\mathbf{x + 1}$
in $\mathrm{x}+\mathrm{b}+\mathrm{a}+1$

## When is a closure created?

let $\mathrm{a}=1 \mathrm{in}$

$$
\text { let } \mathbf{b}=\mathbf{2} \text { in }
$$

$$
\text { let foo }=\text { fun } x->x+1
$$

$$
\text { in } x+b+a+1
$$

## When is a closure created?

## Environment

$$
\begin{aligned}
& \text { let } \mathbf{a}=1 \text { in } \\
& \qquad \begin{array}{l}
\text { let } \mathbf{b}=\mathbf{2} \text { in } \\
\text { let foo }=\text { fun } x \rightarrow x+1 \\
\\
\text { in } \mathbf{x}+\mathrm{b}+\mathbf{a}+1
\end{array}
\end{aligned}
$$

The name is None.
When would it need a name?

## App

## ~Of dark magic~

## How do you call a function?

```
let a = 1 in
let b = 2 in
    let foo = fun X -> X + a
    in foo 5
```


## How do you call a function?

let $\mathbf{a}=1$ in
let $\mathbf{b}=\mathbf{2}$ in
let foo = fun $\mathbf{x}->\mathbf{x}+\mathbf{a}$
in foo 5

## How do you call a function?

let $\mathbf{a}=1$ in
(b,2), (a, 1)
let $\mathbf{b}=\mathbf{2}$ in
let foo = fun $\mathbf{x}$-> $\mathbf{x}+\mathbf{a}$
in foo 5

## How do you call a function?

let $\mathbf{a}=1$ in
(Closure([(b,2),(a,1)], None, $x, x+a)),(b, 2),(a, 1)$
let $\mathbf{b}=\mathbf{2}$ in
let foo = fun $\mathbf{x}$-> $\mathbf{x}+\mathbf{a}$
in foo 5

## How do you call a function?

```
let \(\mathbf{a}=1\) in
let \(\mathbf{b}=\mathbf{2}\) in
    let foo = fun \(\mathbf{x}->\mathbf{x}+\mathbf{a}\)
```

    in foo 5
    
## How do you call a function?

$$
\begin{aligned}
& \text { let } \mathbf{a}=1 \text { in } \\
& \text { let } \mathbf{b}=2 \text { in }
\end{aligned}
$$

$$
\text { let foo = fun } \mathbf{x}->\mathbf{x}+\mathbf{a}
$$

in foo 5
(Closure([(b,2),(a,1)], None, $x, x+a)),(b, 2),(a, 1)$

Assuming we're in eval....

## How do you call a function?

## Environment

let $\mathbf{a}=1$ in<br>let $\mathbf{b}=\mathbf{2}$ in

(Closure([(b,2),(a,1)], None, $x, x+a)),(b, 2),(a, 1)$

Assuming we're in eval....
let foo = fun $\mathbf{x}$-> $\mathbf{x}+\mathbf{a}$
in foo 5

1. We take the environment [(b,2), $(a, 1)]$ inside the closure

## How do you call a function?

## Environment

let $\mathbf{a}=1$ in<br>let $\mathbf{b}=\mathbf{2}$ in

(Closure([(b,2),(a,1)], None, $x, x+a)),(b, 2),(a, 1)$

Assuming we're in eval....
let foo = fun $\mathbf{x}$-> $\mathbf{x}+\mathbf{a}$
in foo 5

1. We take the environment [(b,2), (a,1)] inside the closure
2. Then you bind the parameter $(x, 5)$ to the passed value

## How do you call a function?

## Environment

```
let \(\mathbf{a}=1\) in
let \(\mathbf{b}=\mathbf{2}\) in
```

let foo = fun $\mathbf{x}$-> $\mathbf{x}+\mathbf{a}$
in foo 5

Assuming we're in eval....
(Closure([(b,2),(a,1)], None, $x, x+a)),(b, 2),(a, 1)$

1. We take the environment [(b,2), (a,1)] inside the closure
2. Then you bind the parameter $(x, 5)$ to the passed value
3. Then you pass the new bind [(x,5),(b,2),(a,1)] to the environment

## How do you call a function?

$$
\begin{aligned}
& \text { let } \mathbf{a}=1 \text { in } \\
& \qquad \text { let } \mathbf{b}=\mathbf{2} \text { in }
\end{aligned}
$$

$$
\text { let foo = fun } \mathbf{x}->\mathbf{x}+\mathbf{a}
$$

in foo 5
(Closure([(b,2),(a,1)], None, $x, x+a)),(b, 2),(a, 1)$

Assuming we're in eval....

1. We take the environment [(b,2), (a,1)] inside the closure
2. Then you bind the parameter $(x, 5)$ to the passed value
3. Then you pass the new bind $[(x, 5),(b, 2),(a, 1)]$ to the environment
4. And you evaluate the body in $\begin{aligned} & \text { eval } \\ & {[(x, 5),(b, 2),(a, 1)]}\end{aligned}$ the closure with the new environment

## Whattabout Letrec

Now what happens when you have a recursive function?

## Whattabout Letrec

Now what happens when you have a recursive function?

Whoever gets this right gets a TOBLERONE

## Whattabout Letrec

## Simple.

## Whattabout Letrec

If letrec creates a function, make sure that function has a name!

$$
\text { letrec } \mathbf{x}=\mathbf{e} 1 \text { in } \mathbf{e} 2
$$

## Whattabout Letrec

If letrec creates a function, make sure that function has a name!

## letrec $\mathrm{x}=\mathrm{e} 1$ in e 2

1
Name is 'Some $x$ '

## Map, Fold

## Ghouls, and Ghosts

## To implement the native HOFs

You need to first build the functions using your AST constructors!

## To implement the native HOFs

What are the parameters to map?

## To implement the native HOFs

What are the parameters to map?
('a -> 'b) -> 'a list -> 'b list

## To implement the native HOFs

How do you build a closure for a function named 'map' with takes an argument ' $f$ ?

## To implement the native HOFs

How do you build a closure for a function named 'map' with takes an argument ' $f$ '?

## Closure(env, Some 'map', 'f', <body ....>)

## To implement the native HOFs

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The body is an expression. Thus, its constructed from AST nodes!

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## Closure(env, Some 'map', ‘f', <body ....>)

The body is an expression. Thus, its constructed from AST nodes!

Since we have yet to capture the ' $a$ ' list ' parameter, we may want to start with the Fun constructor. The rest is up to you

## HW4 tips

And other undead creatures

## HW4 in a slide

Problem \#1: evaluate explicit types and binary operations

- Use BinOp's middle argument to find which binary operator (Plus, Minus) is used
- Check for that values have the right type for their operators.
- Else, raise (MLFailure "ERROR TEXT")


## HW4 reminders in a slide

Problem \#1: evaluate explicit types and binary operations

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Problem \#2: Let, Letrec and App

- For the "Lets", you'll be updating the environment. Remember to add the newly named function in the letrec case
- For App, you'll be updating the environment with function parameter


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Problem \#2: Let, Letrec and App

- For the "Lets", you'll be updating the environment. Remember to add the newly named function in the letrec case
- For App, you'll be updating the environment with function parameter

Problem \#3: Native ops

- Start early!

Fin.

## Happy Halloween

