

Name: _____

ID : _____

CSE 130, Fall 2013: Midterm Examination

Nov 5th, 2013

- Do **not** start the exam until you are told to.
- This is a open-book, open-notes exam, but with no computational devices allowed (such as calculators/cellphones/laptops).
- Do **not** look at anyone else's exam. Do **not** talk to anyone but an exam proctor during the exam.
- Write your answers in the space provided.
- Wherever it gives a line limit for your answer, write no more than the specified number of lines. *The rest will be ignored.*
- Work out your solution in blank space or scratch paper, and only put your answer in the answer blank given.
- In all exercises, you are allowed to use the "@" operator.
- Good luck!

1.	20 Points	
2.	20 Points	
3.	30 Points	
TOTAL	70 Points	

1. [20 points] Let's warm up with two small folds. For your reference, the type of `fold_left` is given below:

```
fold_left: ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a
```

a. [10 points] Use `fold_left` to implement `count : 'a list -> 'a -> int`, which returns the number of times a given element occurs in a list. For example:

```
# count [1;2;3;4;5] 10;;
- : int = 0
# count [1;2;3;4;5] 3;;
- : int = 1
# count [1;3;2;3;4;3;5] 3;;
- : int = 3
```

Fill in the implementation of `count` below:

```
let count l x =
```

```
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```

b. [10 points] Use `fold_left` to implement `make_palyndrome : 'a list -> 'a list`, which takes a list and returns a palyndrome, which is produced by adding the elements of the original list in reverse order to the beginning of the original list. For example:

```
# make_palyndrome [1;2];;
- : int list = [2; 1; 1; 2]
# make_palyndrome [1;2;3];;
- : int list = [3; 2; 1; 1; 2; 3]
# make_palyndrome [];;
- : 'a list = []
```

Fill in the implementation of `make_palyndrome` below using `fold_left`:

```
let make_palyndrome l =
```

```
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```

2. [20 points] In this question you will use `fold` to write a slight variant of `fold` and then you will use this variant of `fold` to implement indexing into a list.

a. [10 points] **Fold2.** You will first start by using `fold` to write another function:

```
fold_2 : ('a -> 'b -> int -> 'a) -> 'a -> 'b list -> 'a+
```

This variant of `fold` works exactly like the original `fold`, except that the folding function gets an additional `int` parameter, which is the index of the element that is passed into the folding function. For example the call:

```
fold_2 f "" ["a"; "b"; "c"]
```

would result in

```
(f (f (f "" "a" 0) "b" 1) "c" 2)
```

Fill in the implementation of `fold_2` below using `fold`:

```
let fold_2 f b l =
```


b. [10 points] **Indexing.** You will now use `fold_2` to write a function `ith : 'a list -> int -> 'a -> 'a`, which returns the i^{th} element of a list. In particular, given a list `l`, an integer index `i` greater or equal to 0, and a “default” value `d`, then `(ith l i d)` returns the i^{th} element of the list `l`, or `d` if this element does not exist. For example:

```
# ith ["a";"b";"c";"d"] 0 "";;
- : string = "a"

# ith ["a";"b";"c";"d"] 1 "";;
- : string = "b"

# ith ["a";"b";"c";"d"] 2 "";;
- : string = "c"

# ith ["a";"b";"c";"d"] 3 "";;
- : string = "d"

# ith ["a";"b";"c";"d"] 4 "";;
- : string = ""
```

Fill in the implementation of `ith` below using `fold_2`:

```
let rec ith l i d =
```

```
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```

3. [30 points] Consider the following binary tree datatype:

```
type 'a fun_tree =  
  | Leaf of ('a -> 'a)  
  | Node of ('a fun_tree) * ('a fun_tree);;
```

a. [10 points] **ApplyAll.** You will implement `apply_all : 'a fun_tree -> 'a -> 'a`, which applies all the functions in the tree, using an in-order traversal. For example, suppose we had the following:

```
# let f1 x = x + 1;;  
val f1 : int -> int = <fun>  
  
# let f2 x = x * 2;;  
val f2 : int -> int = <fun>  
  
# let f3 x = x + 3;;  
val f3 : int -> int = <fun>  
  
# let t = Node(Leaf f1, Node(Leaf f2, Leaf f3));;  
val t : int fun_tree = Node (Leaf <fun>, Node (Leaf <fun>, Leaf <fun>))  
  
# apply_all t 0;;  
- : int = 5
```

In particular, `(apply_all t 0)` computes `(f3 (f2 (f1 0)))`. Now fill in the implementation of `fold_2` below using `fold`:

```
let rec apply_all t x =
```

b. [10 points] For each call to `apply_all` below, write down the value returned by `apply_all`:

```
let f1 = (+) 1;;
let f2 = (-) 2;;
let f3 = (+) 3;;
let t = Node(Node(Leaf f1, Leaf f2), Leaf f3);;
```

```
apply_all t 0;;
```

```
let f1 = (^) "a";;
let f2 x = x ^ "b";;
let f3 x = x ^ "ab";;
let t = Node(Leaf f1, Node(Leaf f2, Leaf f3));;
```

```
apply_all t "123";;
```

```
let f1 = List.fold_left (fun x y -> (y*2)::x) [];;
let f2 = List.fold_left (fun x y -> x@[y]) [];;
let t = Node(Node(Leaf f1, Leaf f1), Node(Leaf f1, Leaf f2));;
```

```
apply_all t [1;2;3];;
```

c. [10 points] **Compose.** You will now write a function:

```
compose : 'a fun_tree -> 'a fun_tree -> 'a fun_tree
```

which takes two trees of the same shape and size, and returns a new tree in which the function stored at each leaf is the mathematical composition of the functions stored at the corresponding leaves in the original two trees. Recall that the mathematical composition of two functions f_1 and f_2 is a function $f_3(x) = f_1(f_2(x))$. For example, consider:

```
let t1 = Node(Leaf f1, Leaf f2);;  
let t2 = Node(Leaf f3, Leaf f4);;  
let t3 = compose t1 t2;;
```

In this example, `t3` would be the tree `Node(Leaf f5, Leaf f6)`, where `f5` is the mathematical composition of `f1` and `f3` and `f6` is the mathematical composition of `f2` and `f4`.

Fill in the implementation of `compose` below:

```
let rec compose t1 t2 =
```

```
  match (t1,t2) with
```

```
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```